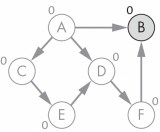
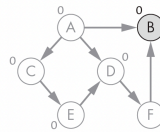


Data Structures the Fun Way

An Amusing Adventure with Coffee-Filled Examples

by Jeremy Kubica

errata updated to print 2

Page	Error	Correction	Print corrected
41	In the next chapter, we will show how can build on these fundamental concepts to create two data structures, stacks and queues, that enable different behavior.	In the next chapter, we will show how we can build on these fundamental concepts to create two data structures, stacks and queues, that enable different behavior.	Pending
169	We start by competing the minimum distance from the point to the node's spatial bounds along each individual dimension:	We start by computing the minimum distance from the point to the node's spatial bounds along each individual dimension:	Pending
200	Formally, we define the size of B-tree node with a size parameter k , which provides bounds on how many elements a non-root node can store.	Formally, we define the size of a B-tree node with a size parameter k , which provides bounds on how many elements a non-root node can store.	Pending
205	We'll explore the later approach, which results in a two-stage algorithm for inserting new keys.	We'll explore the latter approach, which results in a two-stage algorithm for inserting new keys.	Pending
222	B-trees combine indexing and storage in such a way as minimize the number of accesses we need.	B-trees combine indexing and storage in such a way as to minimize the number of accesses we need.	Pending
227	Let's examine how we can extend the hashing techniques we learned Chapter 10 to this prefiltering question.	Let's examine how we can extend the hashing techniques we learned in Chapter 10 to this prefiltering question.	Pending
262	 <p>Next: A Sorted: A,C,E,D,F,B [7]</p> <p>Figure 15-11: A topological sort on a directed acyclic graph</p>	 <p>Next: Sorted: A,C,E,D,F,B [7]</p> <p>Figure 15-11: A topological sort on a directed acyclic graph</p>	Pending